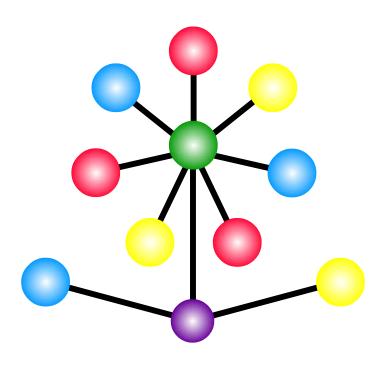
# On strongly almost trivial spatial graphs



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### 1.1 Definitions

G: finite graph

f is a spatial embedding of G

$$\overset{def.}{\Longleftrightarrow} f: G \to R^3: \text{ embedding}$$

We call f (G) a spatial graph

```
f, f': spatial embeddings of G
f and f' are equivalent (f \sim f')
```

$$\stackrel{def.}{\Leftrightarrow}$$
 ∃  $h: R^3 \to R^3$ : orientation preserving self-homeomorphism s.t.  $h \circ f = f'$ 

























### 1.1 Definitions

f is trivial

$$\stackrel{def.}{\Leftrightarrow} \exists f' \sim f$$
s.t.  $f'(G) \subset R^2 \times \{0\} \subset R^3$ 

G is planar

$$\stackrel{def.}{\Leftrightarrow} \exists f: G \rightarrow R^2: \text{ embedding}$$

Hence

G has trivial spatial embeddings  $\iff$  G is planar

























### 1.2 Definitions

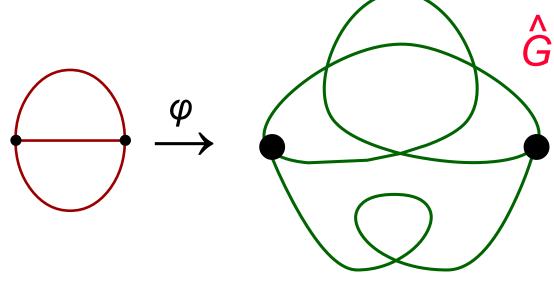
 $\varphi: G \to \mathbb{R}^2$ : continuous map  $\varphi$  is a projection of G

def. multiple points of  $\varphi$  are

only finitely many transversal double points of edges

We call the image of a projection a regular projection

and denote it by  $\hat{G} = \varphi(G)$ 

























### 1.2 Definitions

 $\varphi: G \to \mathbb{R}^2$ : continuous map  $\varphi$  is a projection of G

*def.* multiple points of  $\varphi$  are

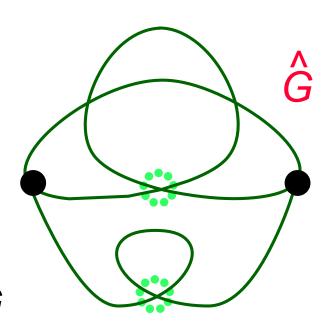
only finitely many transversal double points of edges

We call the image of a projection a regular projection

and denote it by  $\hat{G} = \varphi(G)$ 

A double point of a regular projection is called a crossing

In particular, a crossing point c is a self-crossing def.  $\Leftrightarrow \varphi^{-1}(c) \subseteq e$ , where e is an edge of G



























### 1.2 Definitions

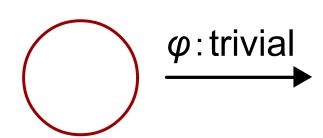
#### $\varphi$ is a projection of a spatial embedding f

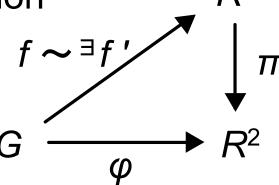
*def.*  $\exists f' \sim f \text{ s.t. } \varphi = \pi \circ f'$  where  $\pi: R^3 \to R^2$  is a natural projection

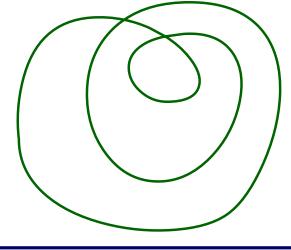
We say f is obtained from  $\varphi$ 

### projection $\varphi$ is trivial

def. only trivial spatial embeddings are obtained from  $\varphi$ 



































### 1.3 Definitions

G: planar graph

f : spatial embedding of G

f is almost trivial.

 $\overset{def.}{\Leftrightarrow} \forall H \subsetneq G$ : proper subgraph,  $f|_H$  is trivial



### 1.3 Definitions

G: planar graph

f : spatial embedding of G

f is minimally knotted.

def. f is nontrivial.

 $\forall H \subsetneq G$ : proper subgraph,  $f|_H$  is trivial

ex.

Brunnian link

























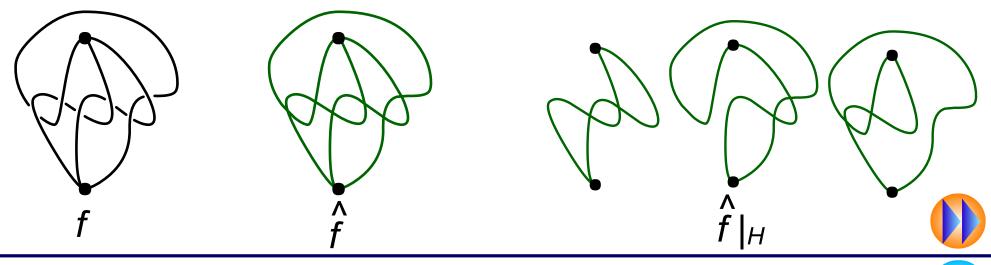


### 1.3 Definitions

G: planar graph, f: spatial embedding of G f is strongly almost trivial (SAT).

 $def. \ f$  is nontrivial,  $\exists \hat{f}$ : projection of f  $\Leftrightarrow$  s.t.  $\forall H \subsetneq G$ : proper subgraph,  $f|_{H}$  is trivial We call  $\hat{f}$  SAT projection.

Hence, f is a SAT embedding of  $G \Rightarrow f$  is minimally knotted ex.  $\theta$ -curve has strongly almost trivial embeddings.

























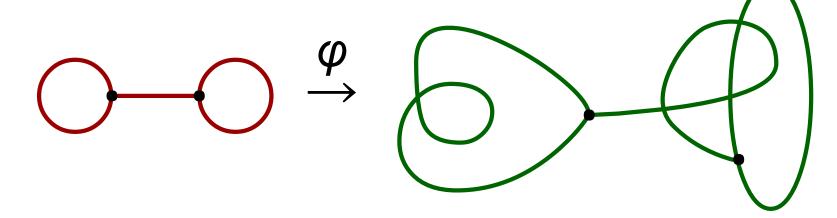


```
\varphi: projection of G \varphi is identifiable.
```

 $\overset{def.}{\Leftrightarrow} \overset{\forall f, f'}{\vdash} : \text{ spatial embeddings from } \varphi$   $f \sim f'$ 

We call  $\varphi$  IP.

ex.





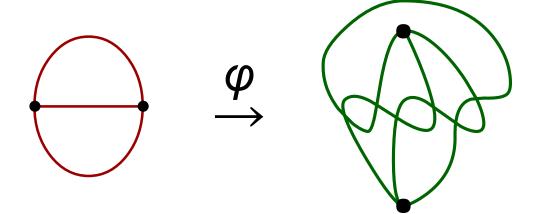
 $\varphi$ : projection of G  $\varphi$  is almost identifiable.

 $\overset{def.}{\Leftrightarrow}$   $\forall H \subsetneq G$ : proper subgraph,  $\varphi|_H$  is identifiable We call  $\varphi$  AIP.

Hence,  $\varphi$  of G is identifiable.  $\Rightarrow \varphi$  is almost identifiable.

 $\varphi$  of G is a SAT projection.  $\Rightarrow \varphi$  is almost identifiable.

ex.





Proposition 1 [Huh-Taniyama, 2002]

Only planar graphs have identifiable projections.

Theorem 2 [Nikkuni, 2005]

 $\varphi$ : projection of G

 $\varphi$  is identifiable

 $\Leftrightarrow f : \text{ spatial embedding obtained from } \varphi,$  f : strivial

























Theorem 3 [Nikkuni, 2005]

G: planar graph which does not have a SAT embedding

 $\varphi$ : projection of G

 $\varphi$  is IP.  $\Leftrightarrow \varphi$  is AIP.

SAT埋め込みをもたないグラフのIPの判定は, 真部分グラフのIPを見ればよい

























Question 1
What kind of planar graphs have minimally knotted spatial embeddings?

Ans. Every planar graph without vertices of degrees ≤ 1 ([Kawauchi, 1989], [Wu, 1993]).



#### Question 2

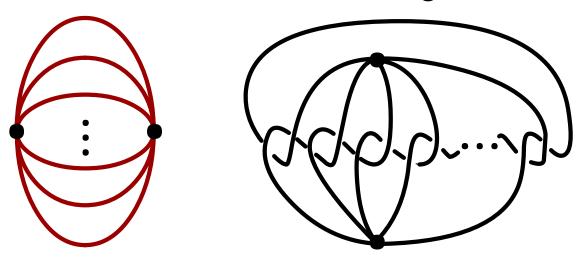
What kind of planar graphs have SAT embeddings?

#### Partial Ans.

<sup>∃</sup>G: planar graph which has a SAT embedding

<sup>∃</sup>G: planar graph which does not have a SAT embedding However it is not well known.

#### ex. $\theta_n$ -curve has SAT embeddings















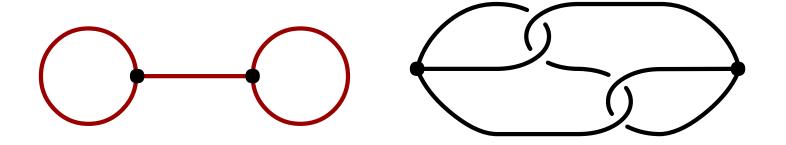








ex. Handcuff graph has SAT embeddings.























Theorem 4 [Huh-Oh, 2003]

G: connected planar graph which does not have a cut vertex G satisfies the followings

- (1) G has no multiple edges.
- (2)  $\forall e_1, e_2 \in E(G)$  s.t.  $e_1 \cap e_2 = \emptyset$ ,  $\exists C_1, C_2$ : disjoint cycles s.t.  $e_1 \in E(C_1)$ ,  $e_2 \in E(C_2)$

$$e_1$$
  $C_1$   $e_2$   $C_2$ 

(3)  $\forall e_1, e_2, e_3 \in E(G)$  s.t.  $e_1 \cup e_2 \cup e_3$  is homeo. to a path,  $\exists C$ : cycle s.t.  $e_1, e_2, e_3 \in E(C)$   $e_2 \cap C$ 

 $\Rightarrow$  G has no strongly almost trivial embeddings.















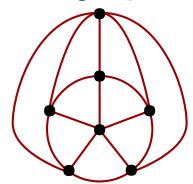




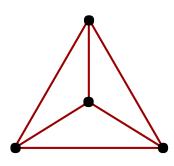




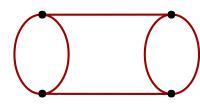
ex. graphs which does not have a SAT embedding



P<sub>5</sub> satisfies all assumptions of Thm.4



*K*<sub>4</sub> does not satisfy the assumption (2) of Thm. 4, but it does not have a SAT embedding [Huh-Oh, 2002]



Double-handcuff graph does not satisfy the assumptions (1) and (2), but it does not have a SAT embedding [Hanaki, preprint]























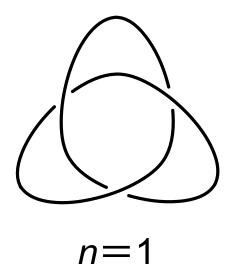


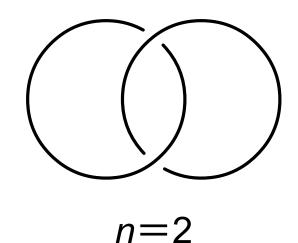


Question 3 Does  $G=S_1^1 \perp S_2^1 \perp \cdots \perp S_n^1$  have SAT embeddings?

Ans. It has SAT embeddings if n = 1,2It has no SAT embeddings if  $n \ge 3$ 

ex.



























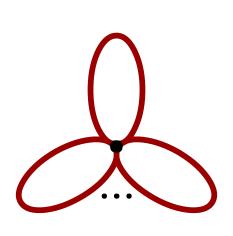


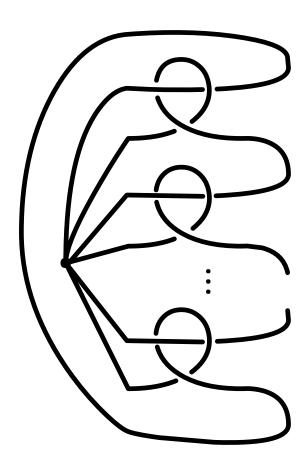


# 3.1 わかったこと

#### Main Theorem 1 [Hanaki]

*n*-bouquet has SAT embeddings.



















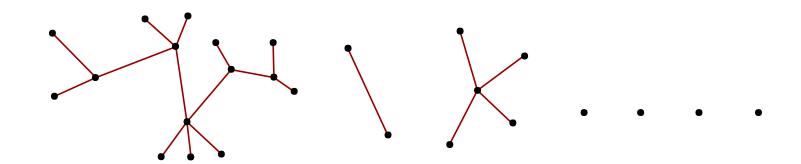




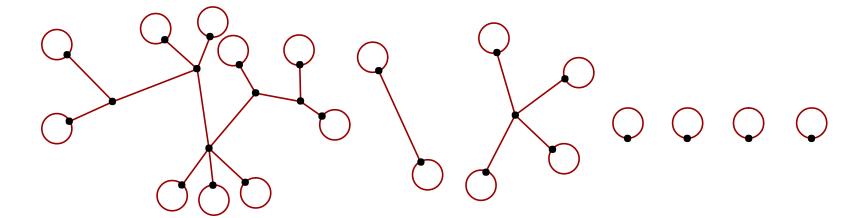




F: forest



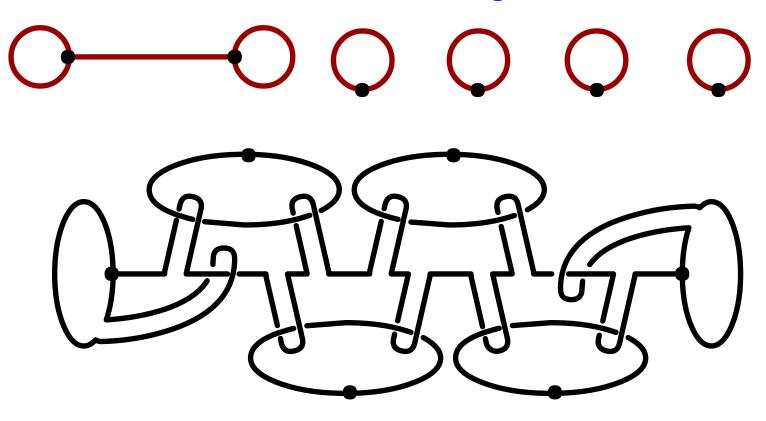
*G<sub>F</sub>*: graph obtained from *F* 





#### Main Theorem 2 [Hanaki]

If  $E(F) \neq \emptyset$ ,  $G_F$  has SAT embeddings.



















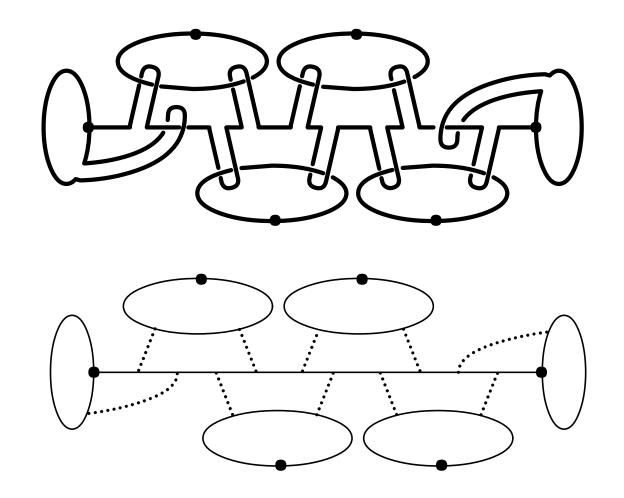








#### cord presentation





















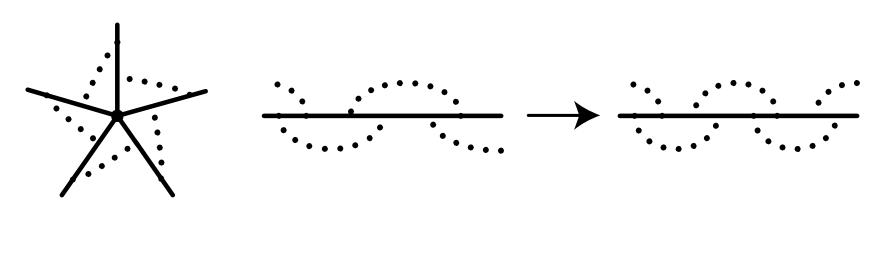


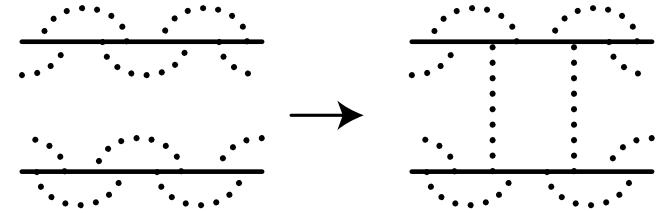






#### SATの構成方法





















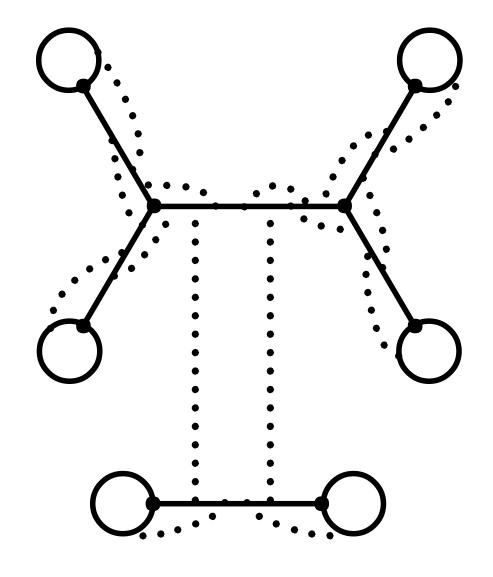




































# 3.3 わかったこと3

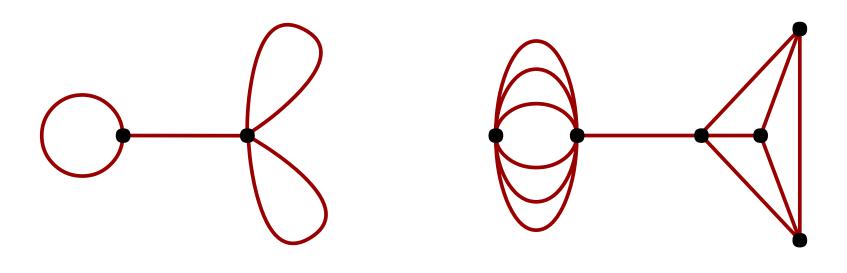
#### Main Theorem 3 [Hanaki]

G: conn. graph s.t. G is not homeo. to handcuff graph  $\exists e \in E(G)$  s.t. e is a cut edge

 $H_1$ ,  $H_2$ : conn. comp. of G-e

 $H_1$ ,  $H_2$  has no cut edges and has cycles

 $\Rightarrow$  G has no SAT embeddings























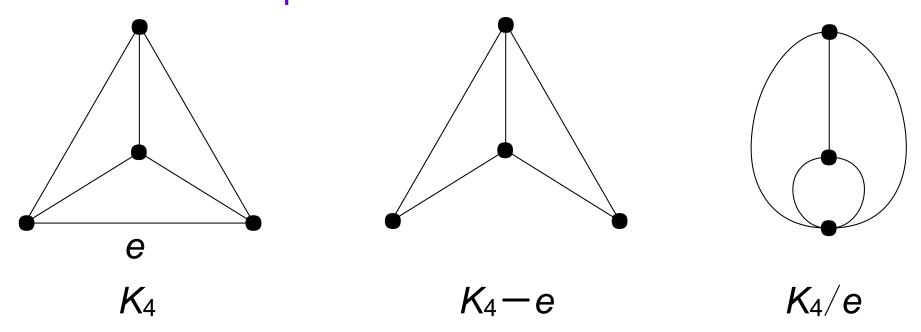






Edge deletion (G-e) is the process of removing only an edge.

Edge contraction (G/e) is the process of removing an edge and combining its two endvertices into a single vertex where e is not loop





*H* is a minor of G ( $H <_m G$ )

def. H can be obtained from G by contracting edges, deleting edges, and deleting isolated vertices

A property  $\mathcal{P}$  of a graph is inherited by minors

 $\overset{def.}{\Leftrightarrow}$  If G has  $\mathcal{P}$ ,  $\forall H <_m G$ , H has  $\mathcal{P}$ 

#### Robertson-Seymour's Minor Theorem

G has  $\mathcal{P}$ 

 $\Leftrightarrow$  G does not contain  $G_1, G_2, \cdots$  and  $G_n$  as a minor

finite





















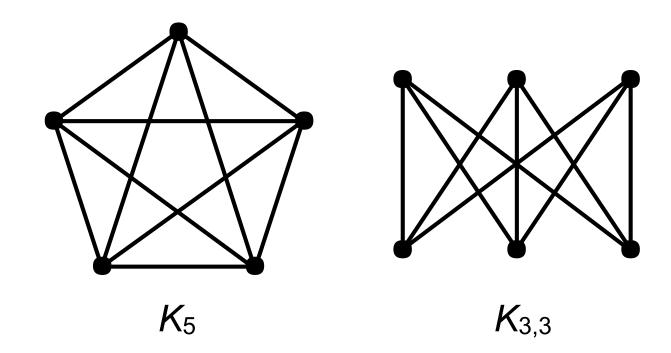




ex. Kuratowski Theorem

G is planar

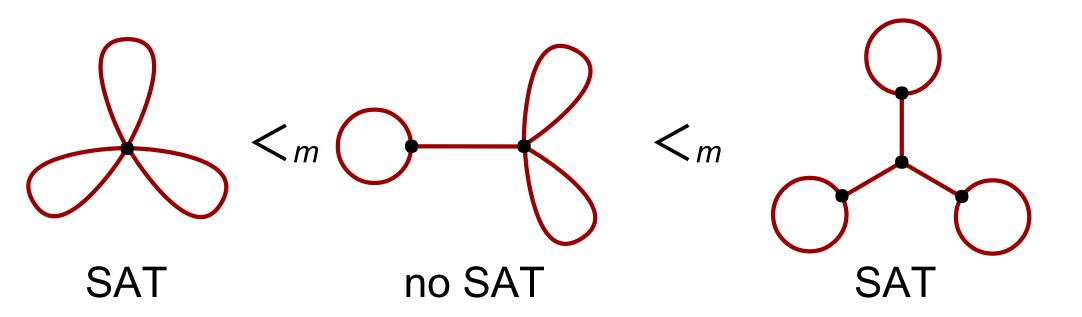
 $\Leftrightarrow$  G does not contain  $K_5$ ,  $K_{3,3}$  as a minor





#### Remark

A property that a graph has (no) SAT embeddings is not inherited by minors.



























 $G \subset S^3$  is irreducible

 $G \subset S^3$ : spatial graph

 $D \subset S^3$ : disk

D is good for G

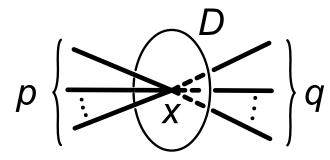
*def.*  $\partial D \subseteq G$ 

 $intD \cap G$  contains at most finitely many points

 $x \in \text{int}D \cap G$ ,

a neightbourhood of x looks like

where  $p, q \in N$ 

























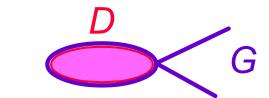




Theorem [Taniyama, 2002]

 $G \subset S^3$ : spatial graph

 $D \subset S^3$ : disk s.t. D is good for G



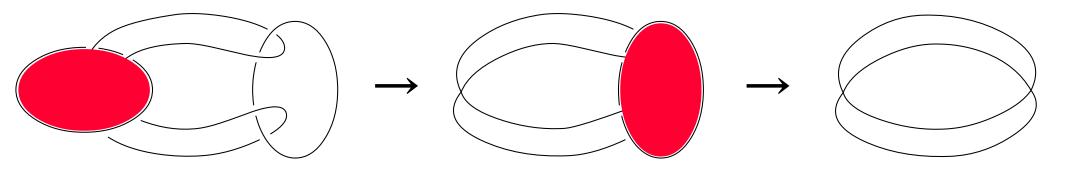
int $D \cap G = \emptyset$  or  $\partial D \cap \operatorname{cl}(G - \partial D)$  is not singleton where cl denotes the closure

 $G' \subset S^3$ : spatial graph obtained from G by contracting D to a point

G' is irreducible  $\Rightarrow$  G is irreducible



ex.





ex.

